

Algorithm Problem Solving (APS): Dynamic Programming

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UC San Diego SPIS 2019

Example: The 0-1 Knapsack Problem

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- Imagine you're in a heist crew en route to your next job
- Your target is a truck carrying electronic goods
- The modified Honda Civic you're driving has a weight limit
 - You know how much each item costs and weighs
 - Which items do you steal to maximize your profits?

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- **Output:** A set of items such that the sum of their weights is below x and the sum of their values is **maximized**

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$w:$	3	8	10	6
$v:$	\$5	\$5	\$13	\$6

$$x = 20$$

Ex

Can we use a Greedy Algorithm?

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$$x = 20$$

Nope! Not guaranteed to be optimal!

Simpler Example: Fibonacci Sequence

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- The first number is 0

0

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- The first number is 0
- The second number is **1**

0, **1**

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0, 1, 1, 2, **3**

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0, 1, 1, 2, 3, 5

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0, 1, 1, 2, **3**, **5**

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0, 1, 1, 2, 3, 5, **8**

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0, 1, 1, 2, 3, 5, 8

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0, 1, 1, 2, 3, 5, 8, **13**

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0, 1, 1, 2, 3, 5, 8, 13, ...

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Simpler Example: Fibonacci Sequence

- Write a recursive Python function!
- **Output:** The n -th number of the Fibonacci sequence (0-based indexing)

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Recursive Fibonacci

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def fib_rec(n):  
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        return 0  
  
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Hold up... 🤔

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rec_fib(5)

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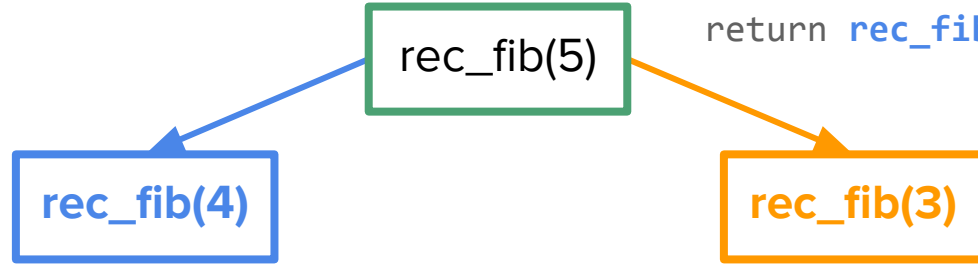
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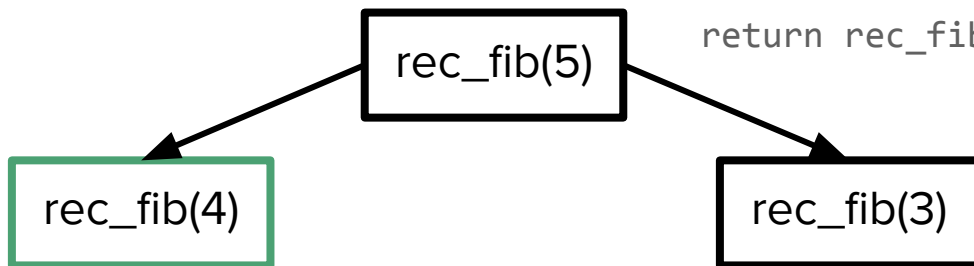
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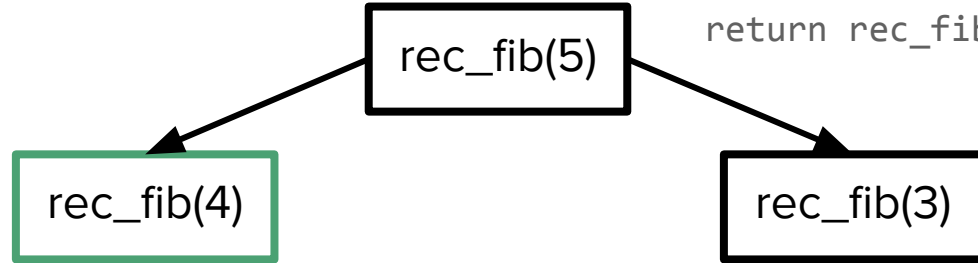
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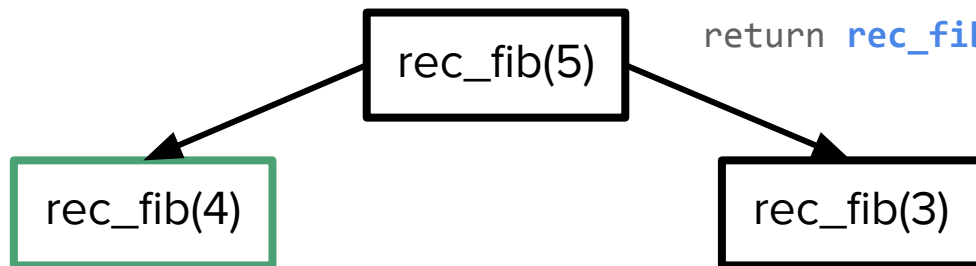
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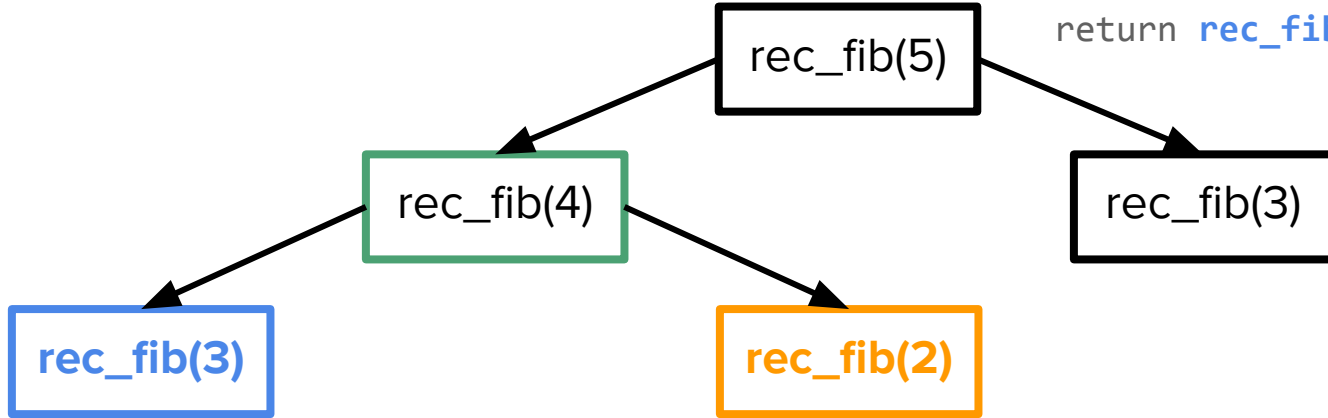
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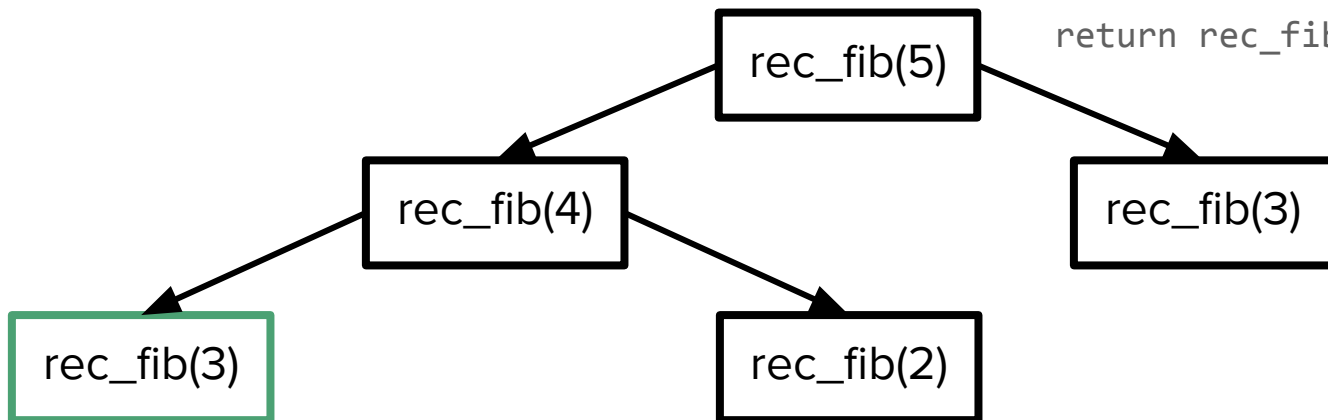
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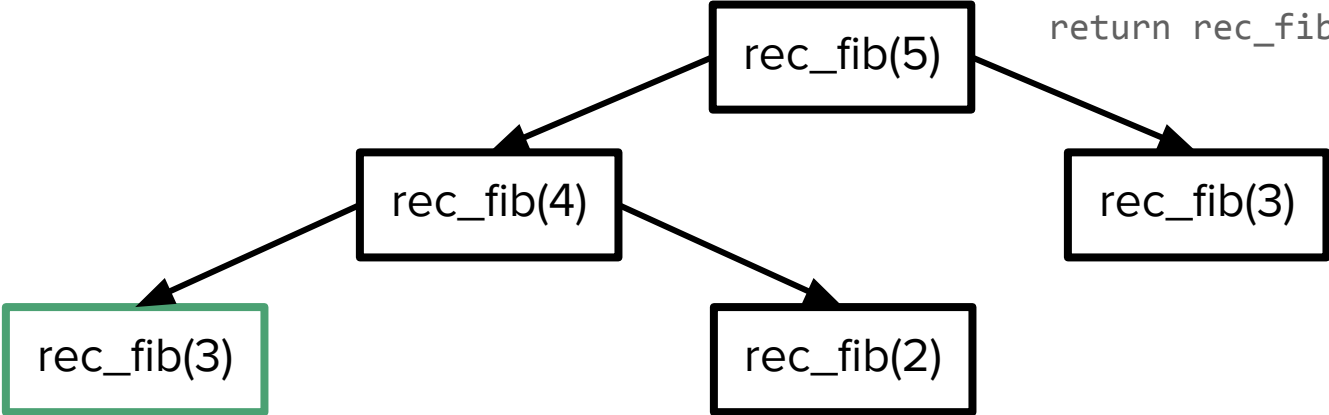
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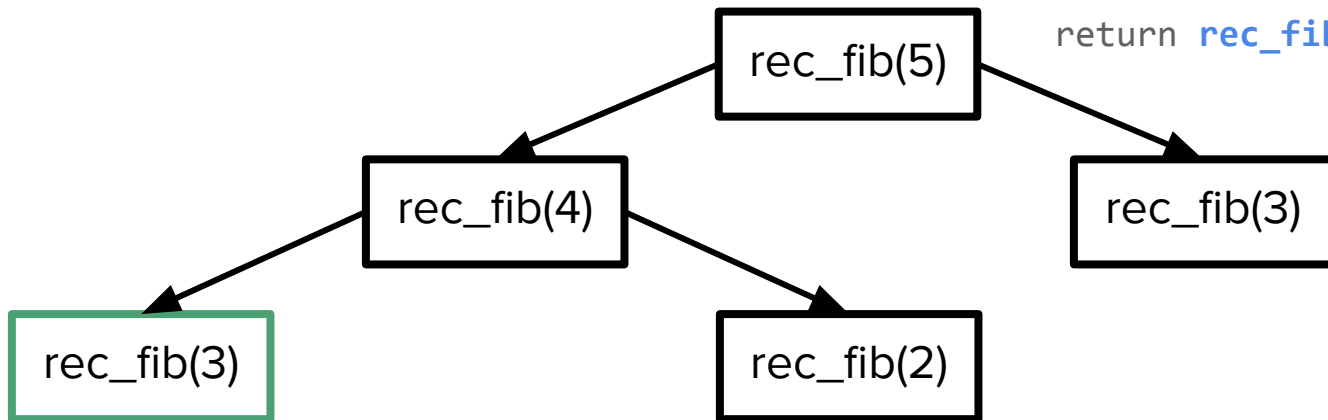
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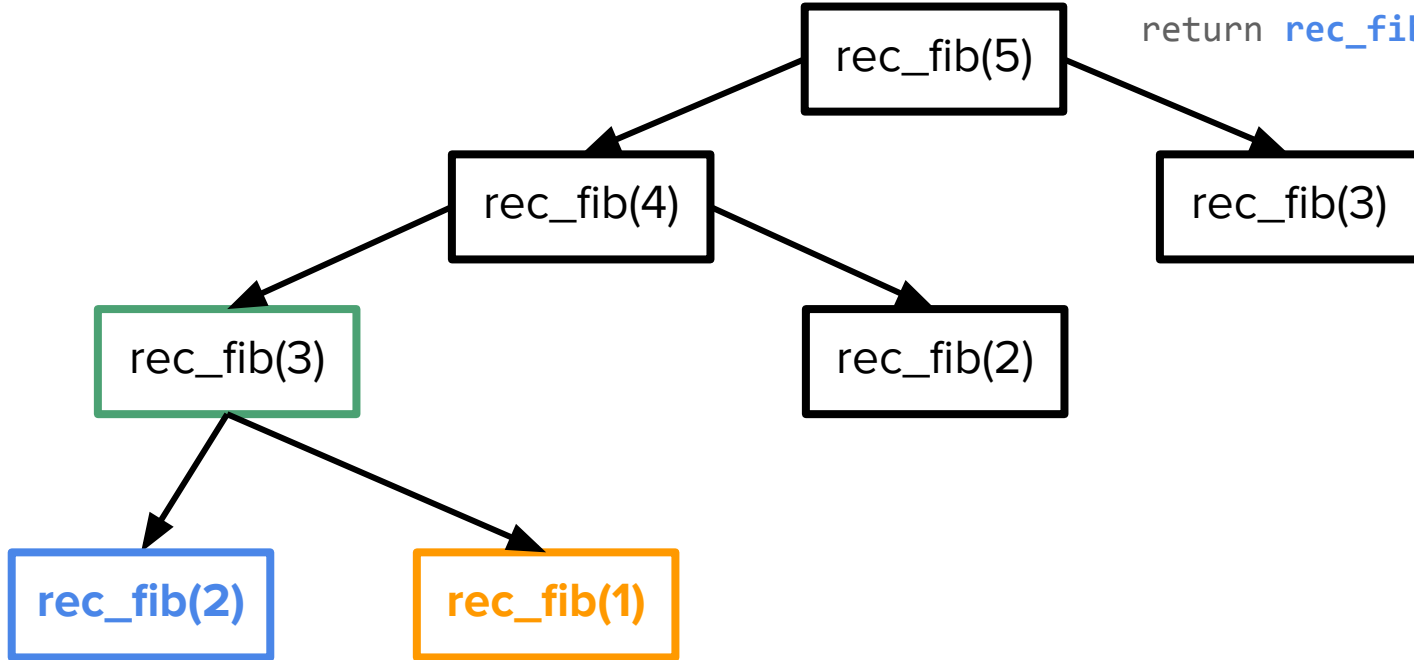
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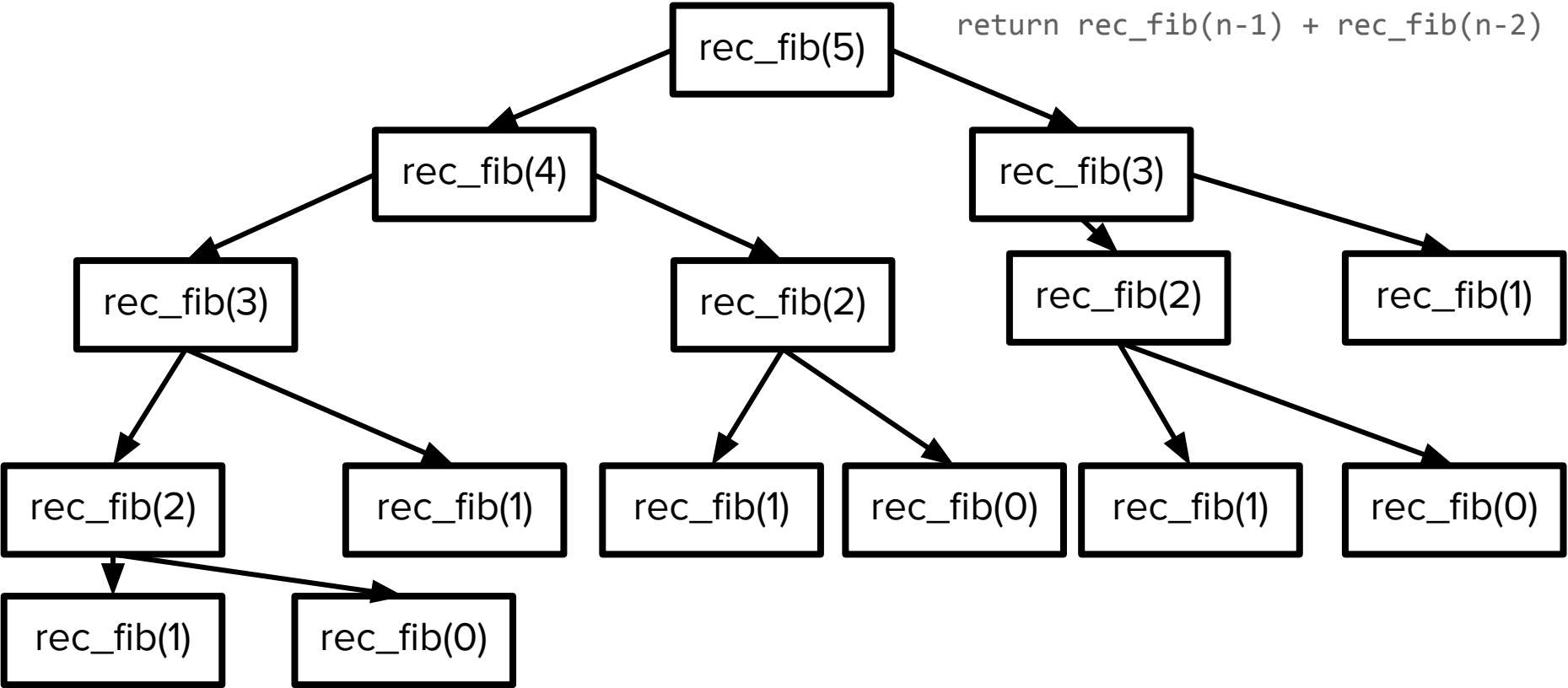
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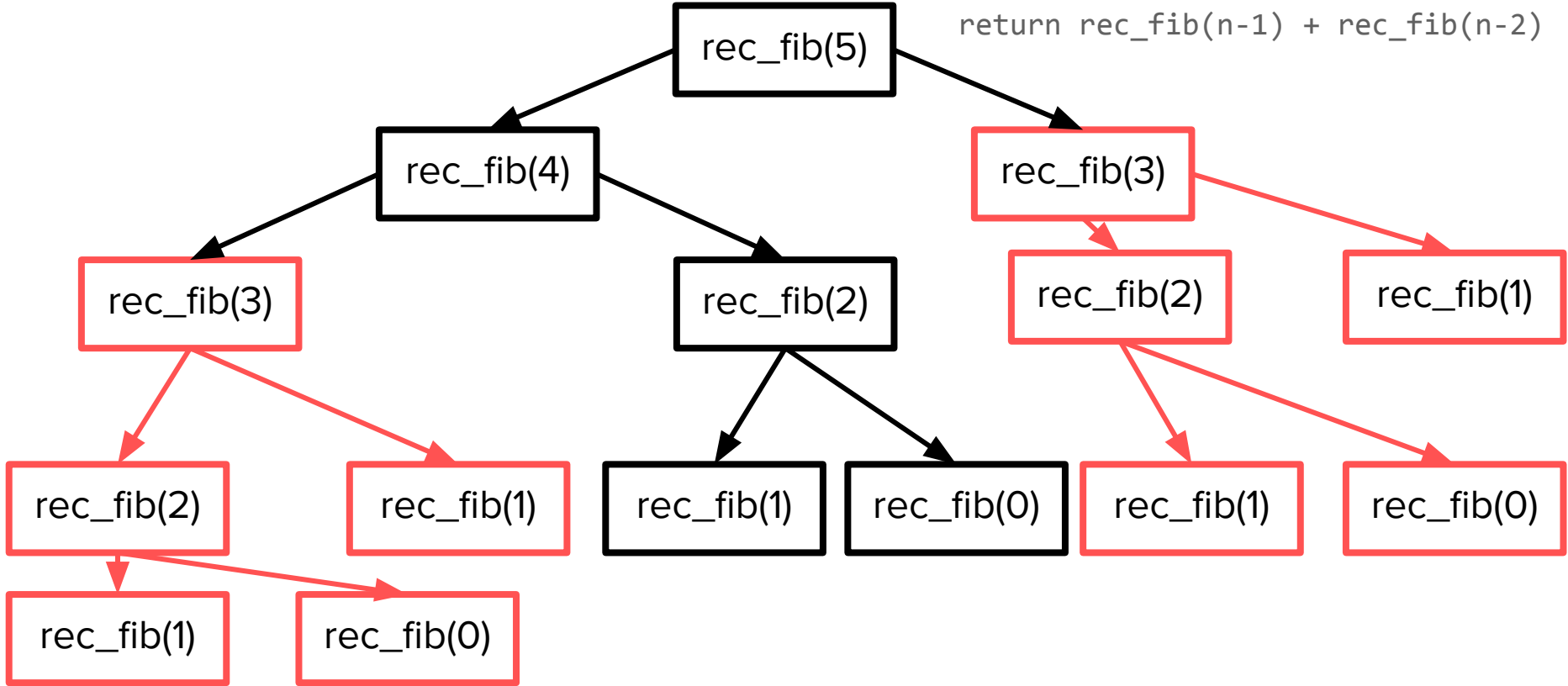
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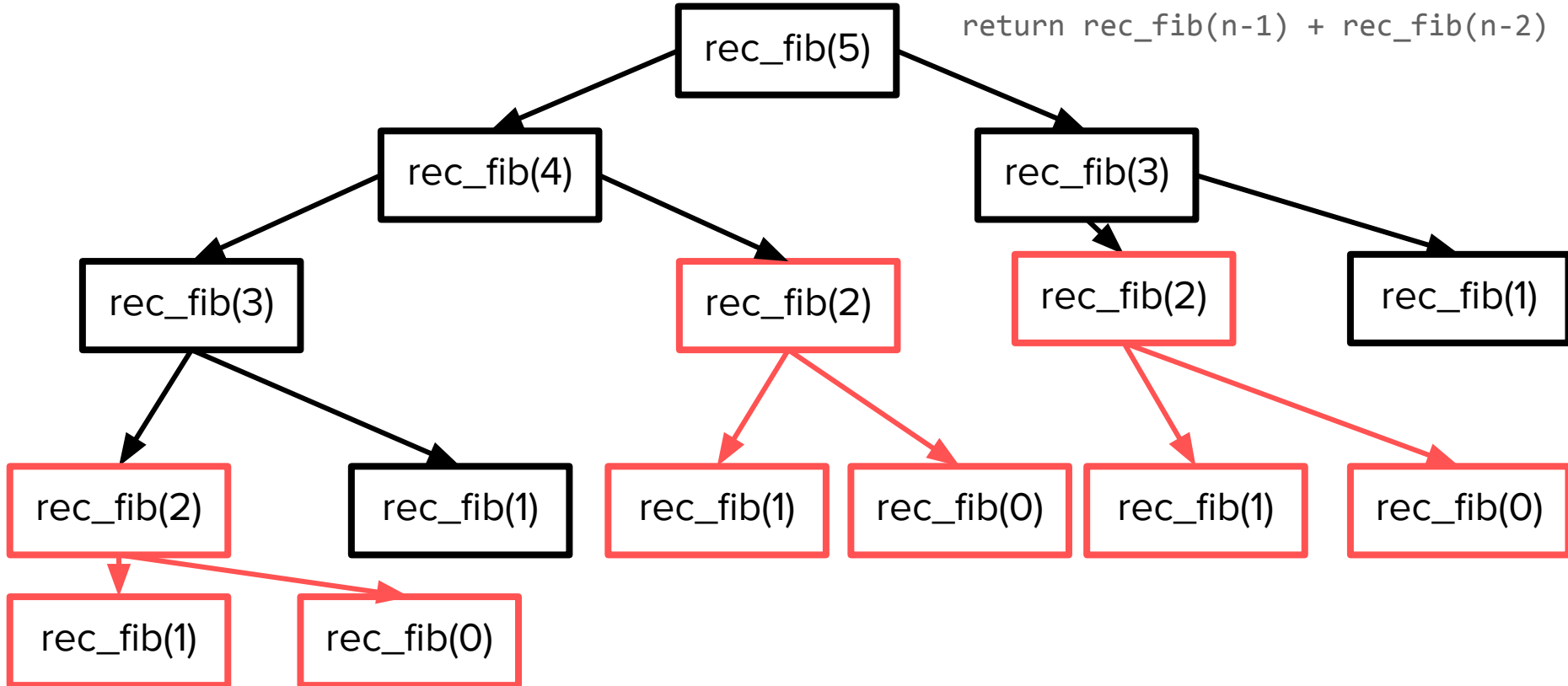
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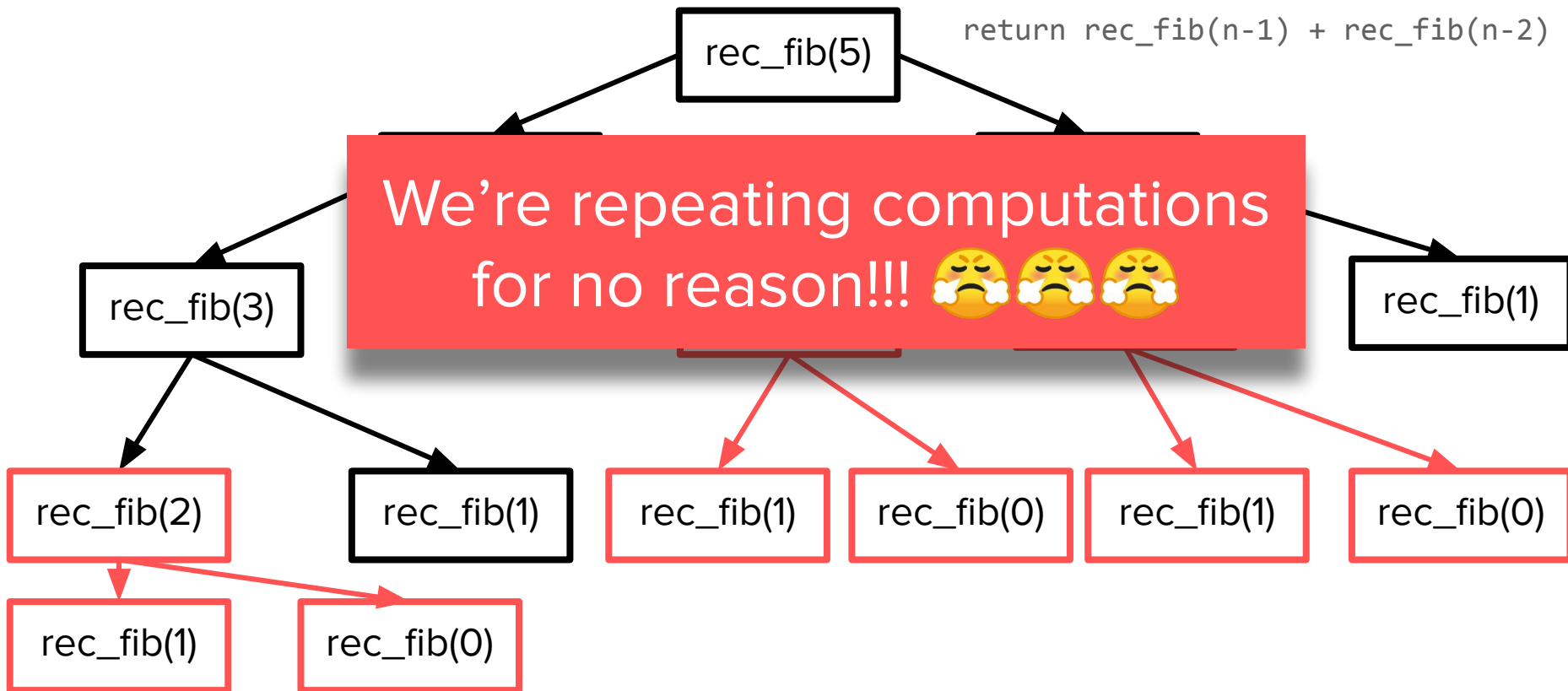
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 - The values of `rec_fib(3)` and `rec_fib(2)` never change, so once we compute them the first time, why not just save them somewhere?
- **Memoization:** Saving the results of expensive functions somewhere and returning the saved result when the same inputs occur again

Fast Fib(e)

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fast_fib(5)
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fast_fib(5)

<i>n</i> :	0	1	2	3	4	5
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 - Don't add item i to our pack, so our profit would be $P(i-1, j)$

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 - Don't add item i to our pack, so our profit would be $P(i-1, j)$
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 - Don't add item i to our pack, so our profit would be $P(i-1, j)$
 - No value added, so it's just our best scenario of looking at the first $i-1$ items with the same (j) total weight
 - Take the **maximum** of these two possible options

Example: The 0-1 Knapsack Problem

$$P(i, j) = \begin{cases} 0 & \text{if } i = 0 \\ -\infty & \text{if } j < 0 \\ \max \{v_i + P(i - 1, j - w_i), P(i - 1, j)\} & \text{if } 1 \leq i \leq n, 0 \leq j \leq x \end{cases}$$

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v :	\$5	\$5	\$13	\$6

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j (total weight)

i (number of items seen)	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0									0	0	0	0	0	0
1	?																				
2																					
3																					
4																					

Add: \$5 + $-\infty$

Example: The 0-1 Knapsack Problem

$x = 20$

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$$P(i, j) = \begin{cases} 0 & \text{if } i = 0 \\ -\infty & \text{if } j < 0 \\ \max \{v_i + P(i - 1, j - w_i), P(i - 1, j)\} & \text{if } 1 \leq i \leq n, 0 \leq j \leq x \end{cases}$$

j (total weight)

i (number of items seen)	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0									0	0	0	0	0	0
1	?																				
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3																					
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Add: \$5 + $-\infty$

Don't Add: \$0

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i (number of items seen)	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	
0	0	0	0	0	0	0	0									0	0	0	0	0	0	
1	0																					
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	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	5																	
2	0	0	0	5																	
3	0	0	0																		
4	0	0	0																		

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0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	5																		
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j (total weight)

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	5	5																
2	0	0	0	5	5																
3	0	0	0	5	5																
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j (total weight)

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	5	5	5															
2	0	0	0	5	5	5															
3	0	0	0	5	5	5															
4	0	0	0	5	5	5															

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	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0	0
1	0	0	0	5	5	5	5														
2	0	0	0	5	5	5	5														
3	0	0	0	5	5	5	5														
4	0	0	0	5	5	5															

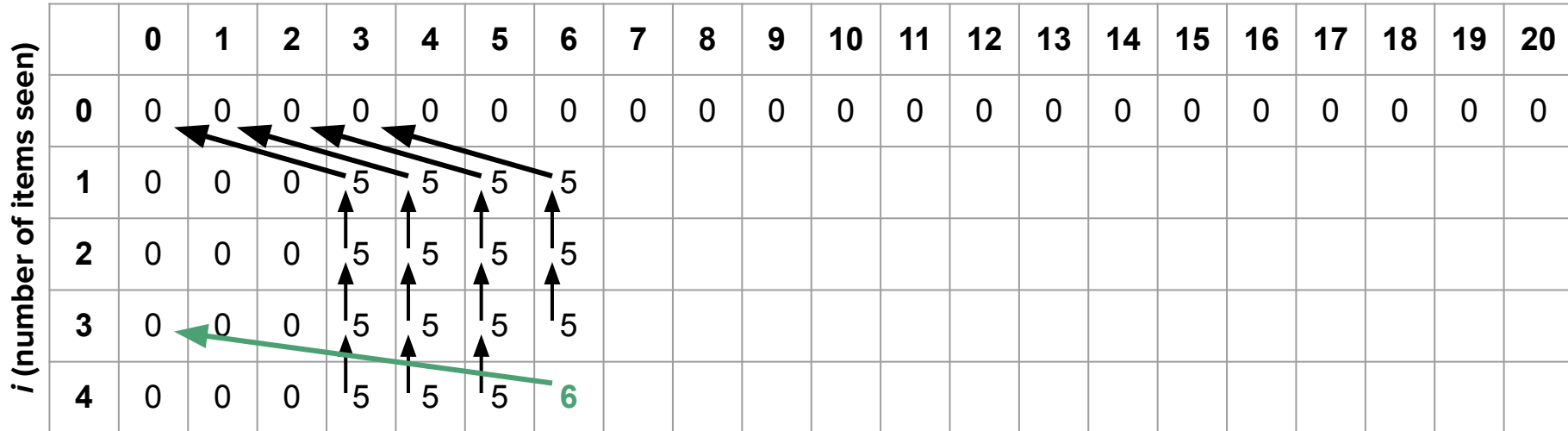
Example: The 0-1 Knapsack Problem

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$w:$	3	8	10	6
$v:$	\$5	\$5	\$13	\$6

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j (total weight)



Example: The 0-1 Knapsack Problem

$x = 20$

w :	3	8	10	6
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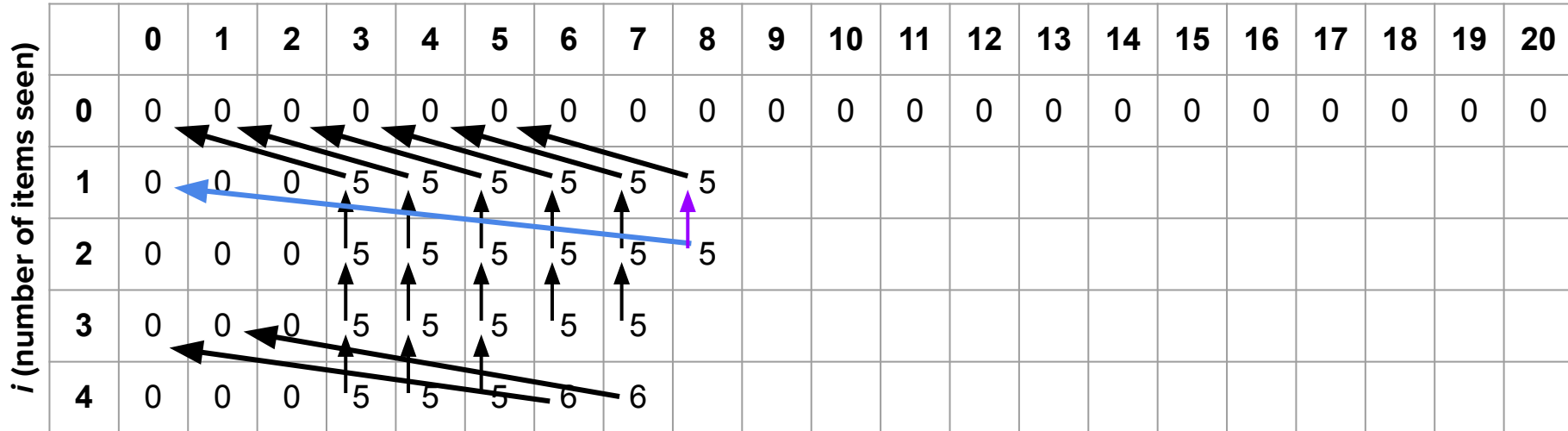
$$P(i, j) = \begin{cases} 0 & \text{if } i = 0 \\ -\infty & \text{if } j < 0 \\ \max \{ v_i + P(i-1, j-w_i), P(i-1, j) \} & \text{if } 1 \leq i \leq n, 0 \leq j \leq x \end{cases}$$

j (total weight)

if $i = 0$

if $j < 0$

if $1 \leq i \leq n, 0 \leq j \leq x$



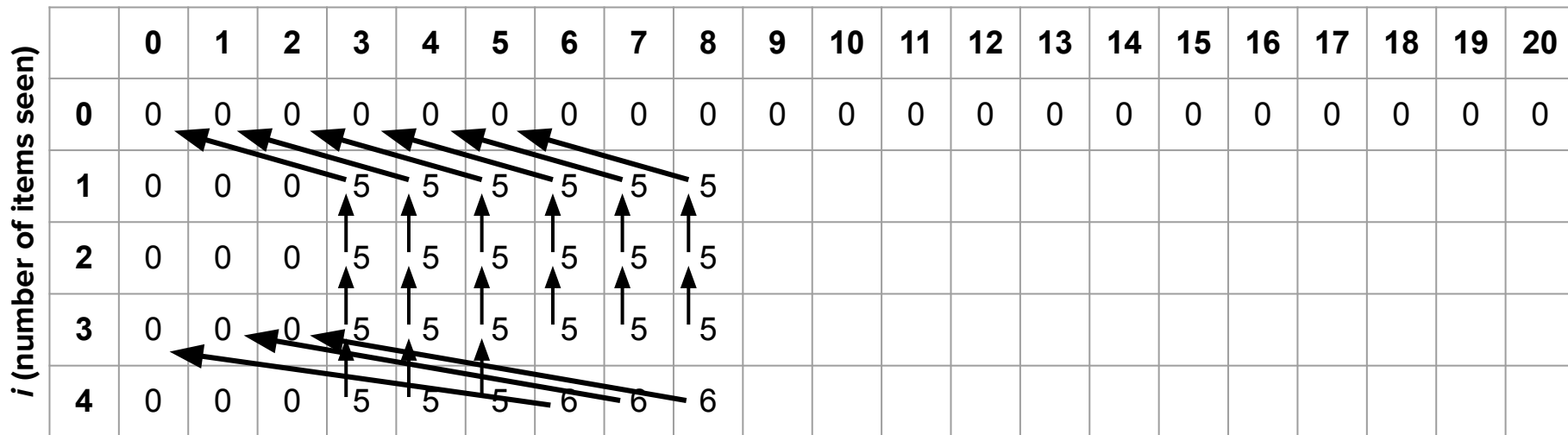
Example: The 0-1 Knapsack Problem

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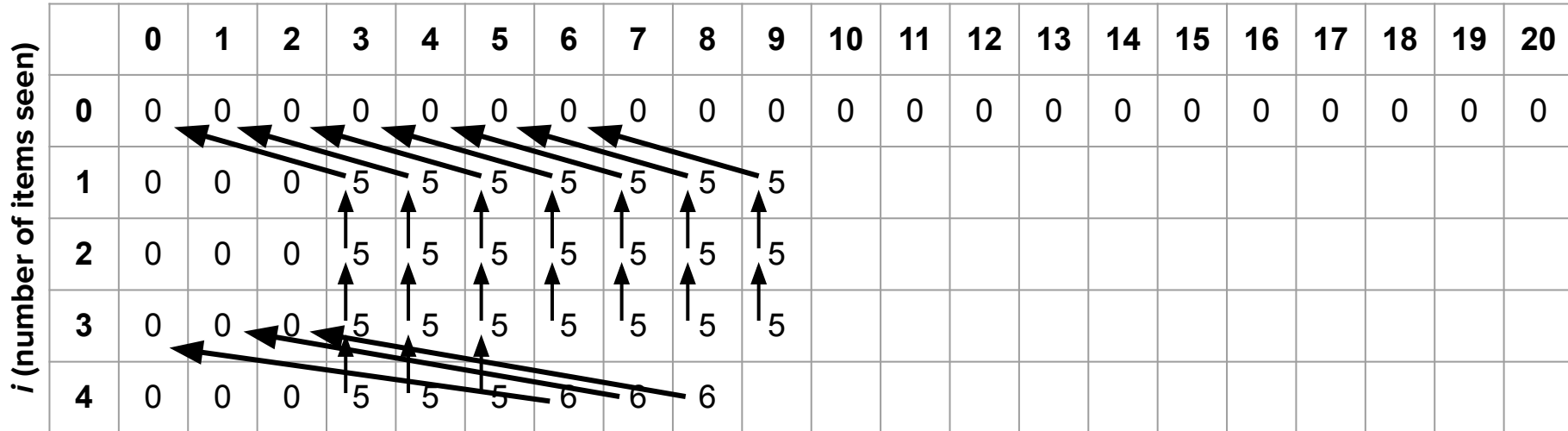
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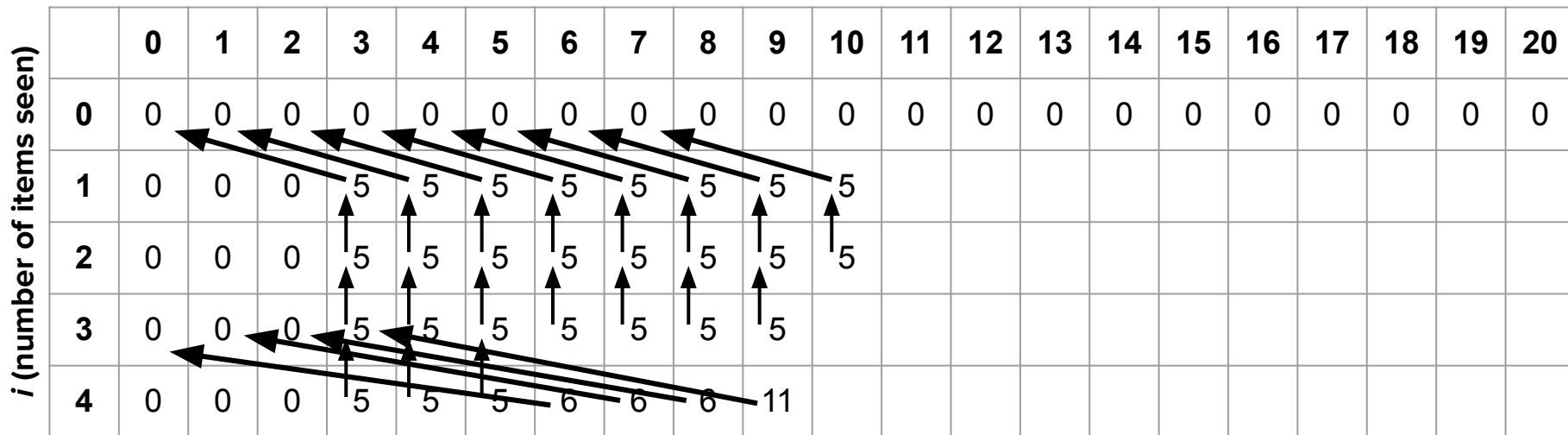
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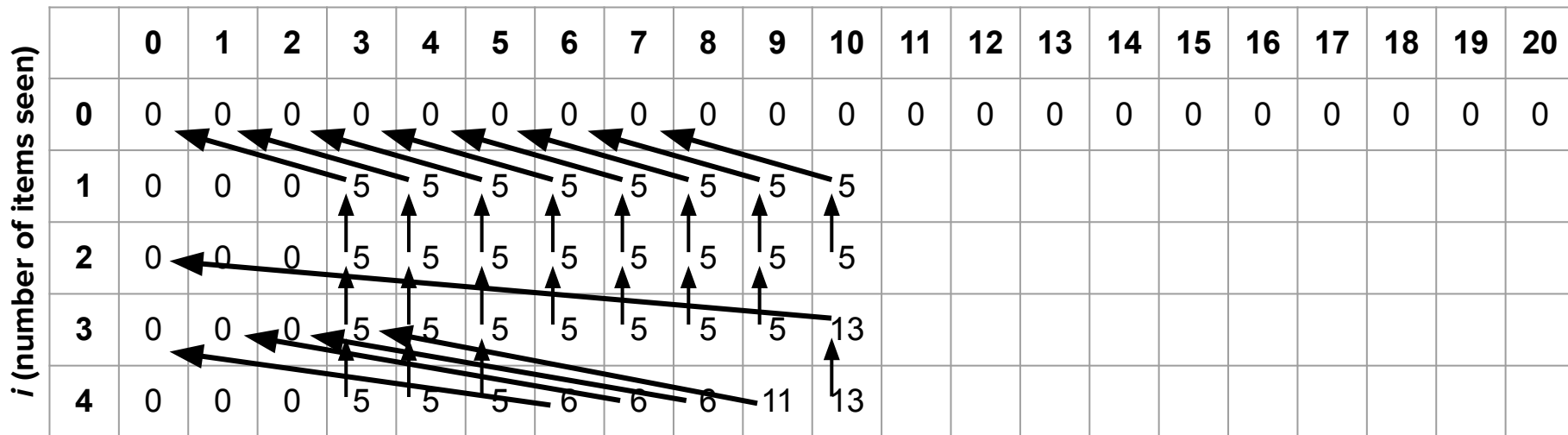
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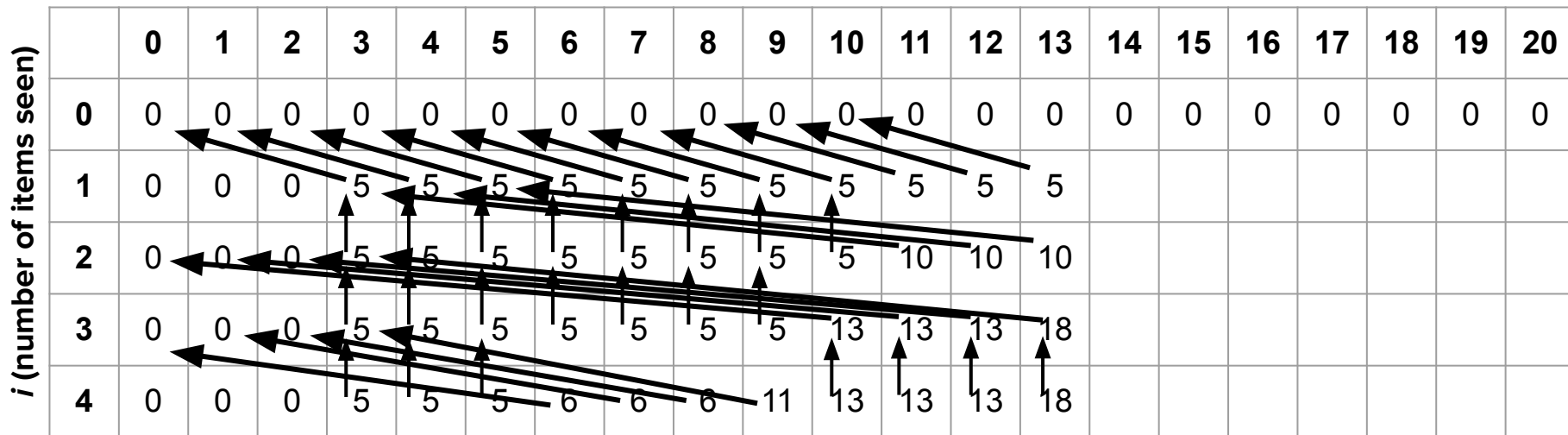
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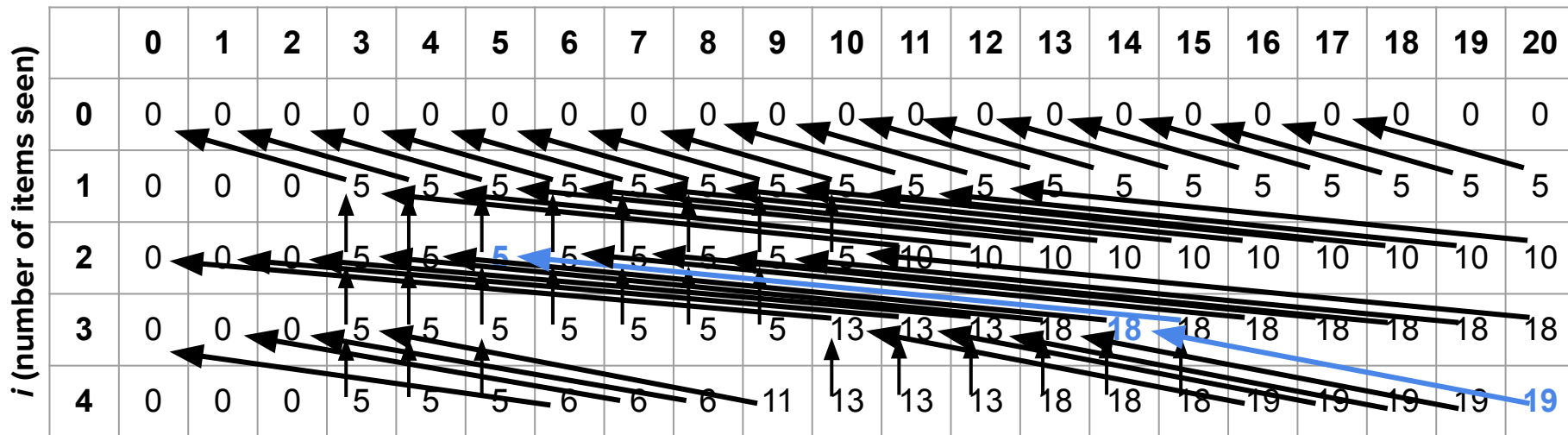
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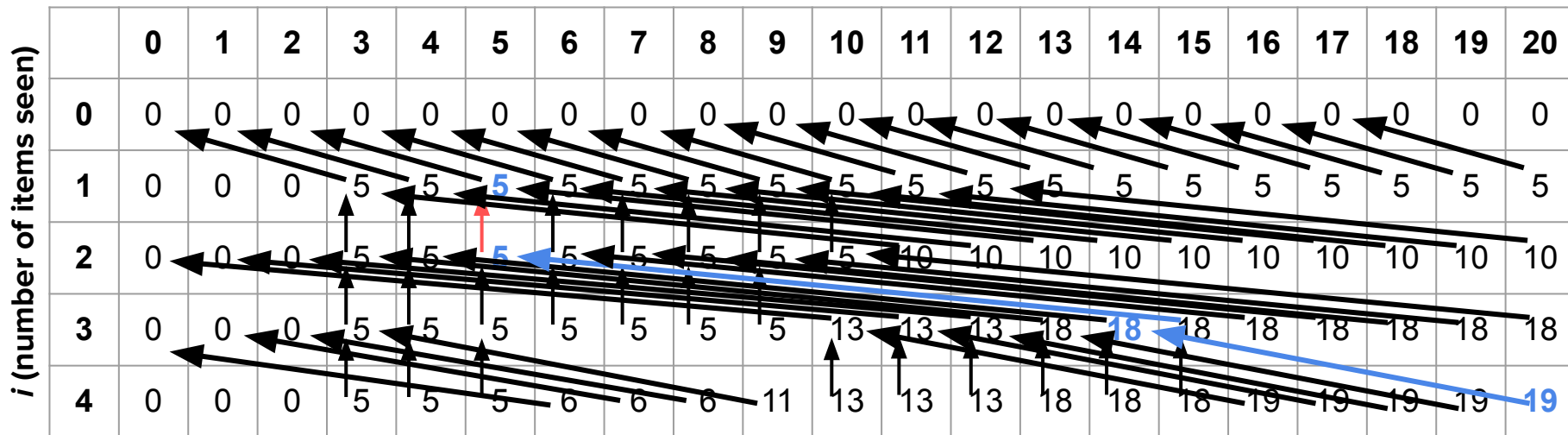
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